

EFFICIENT DATA STRUCTURE
FOR
REPRESENTING & SIMPLIFYING
SIMPLICIAL COMPLEXES
IN
HIGH DIMENSIONS

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MOTIVATION

Reconstruction

Simplification

Analysis



3D scans

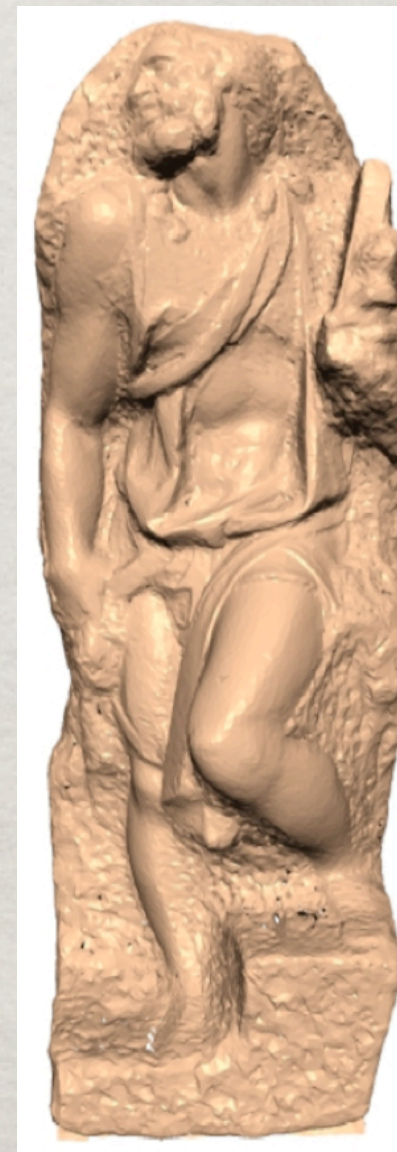
large
data sets



$\approx 400,000,000$ triangles

gigantic
simplicial complex

21

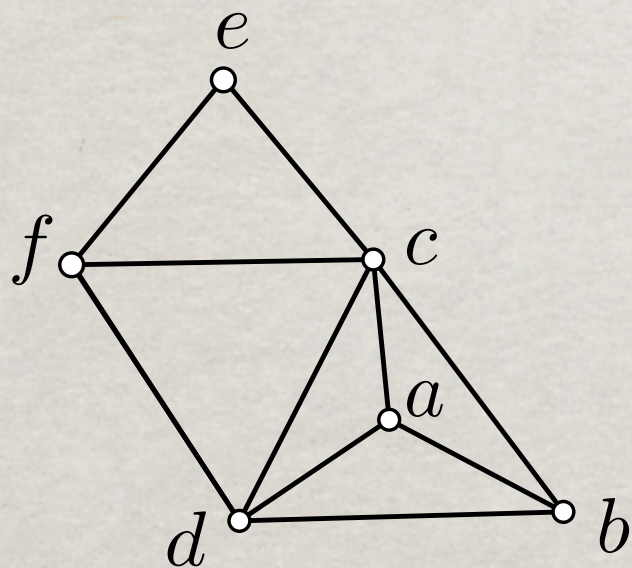


$\approx 100,000$ triangles

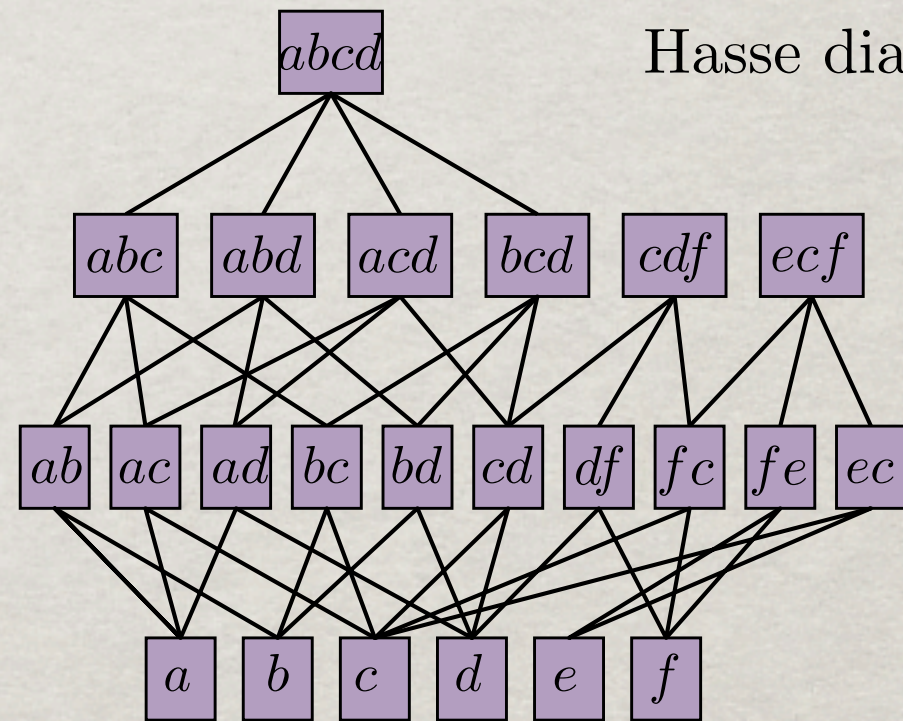
reduced
simplicial complex

FLAG COMPLEXES

- ✱ Flag $G =$ largest simplicial complex whose 1-skeleton is G .
- ✱ $\{v_0, \dots, v_k\} \in \text{Flag } G \iff v_i v_j \in G \ \forall i, j$



G



Hasse diagram

Flag G

Flag complexes have a compressed form of storage

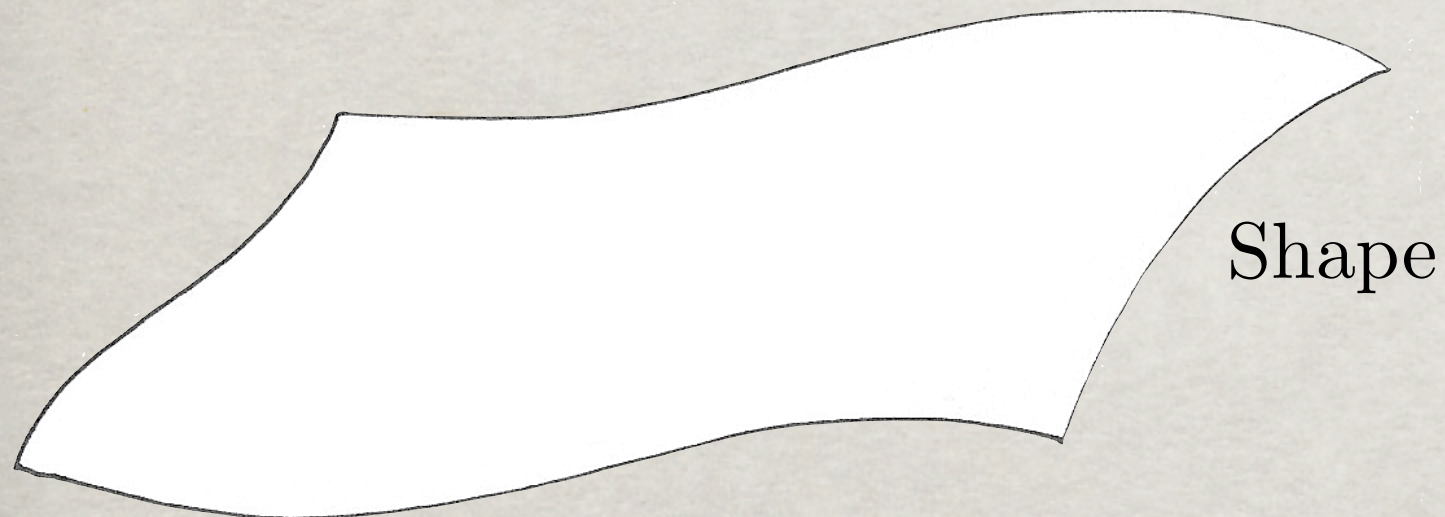
SHAPE RECONSTRUCTION

INPUT

Point cloud P

OUTPUT

Flag $G_\alpha(P) = \text{Rips complex}$



Shape

$G_\alpha(P) = \text{proximity graph}$

$$pq \in G_\alpha(P) \iff d(p, q) \leq 2\alpha$$

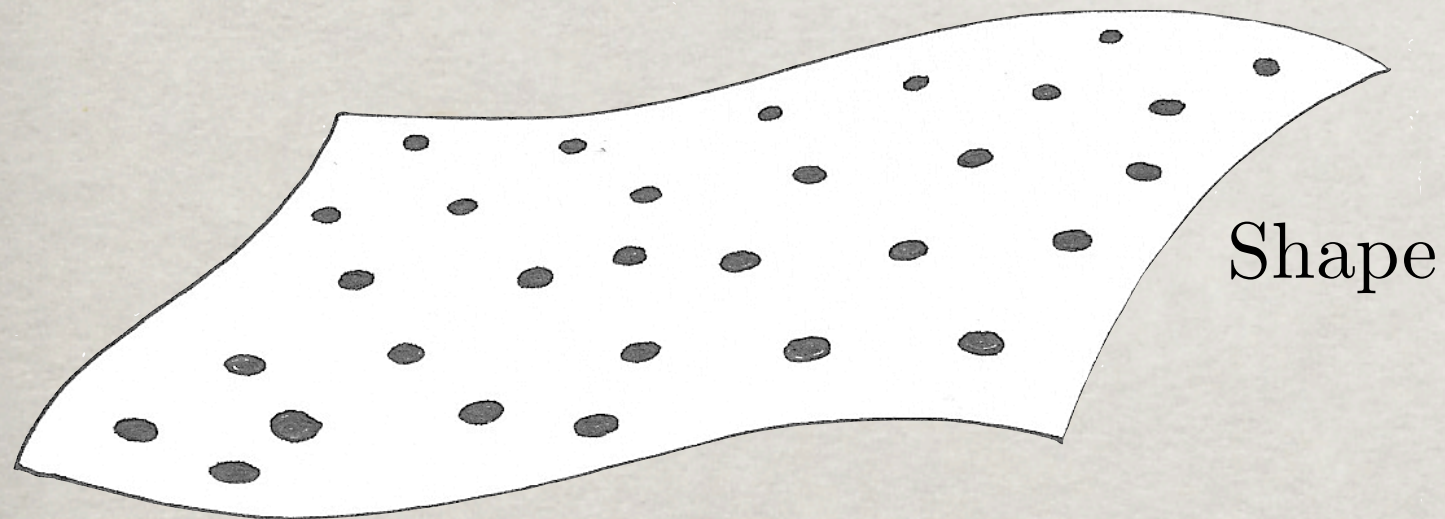
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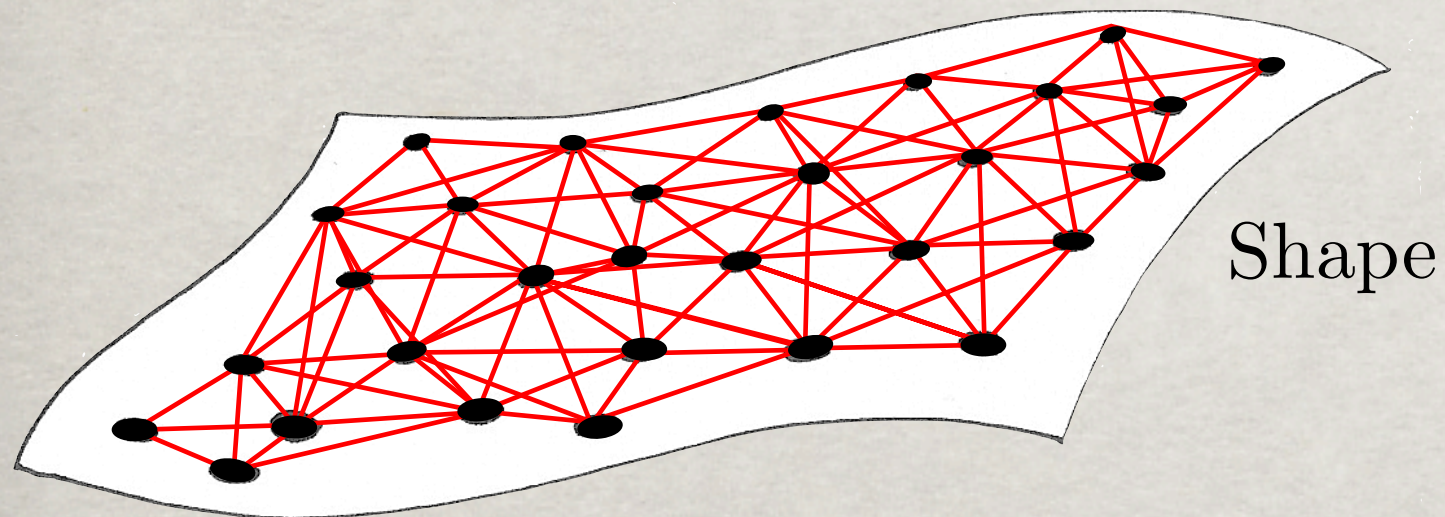
SHAPE RECONSTRUCTION

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Flag $G_\alpha(P) = \text{Rips complex}$



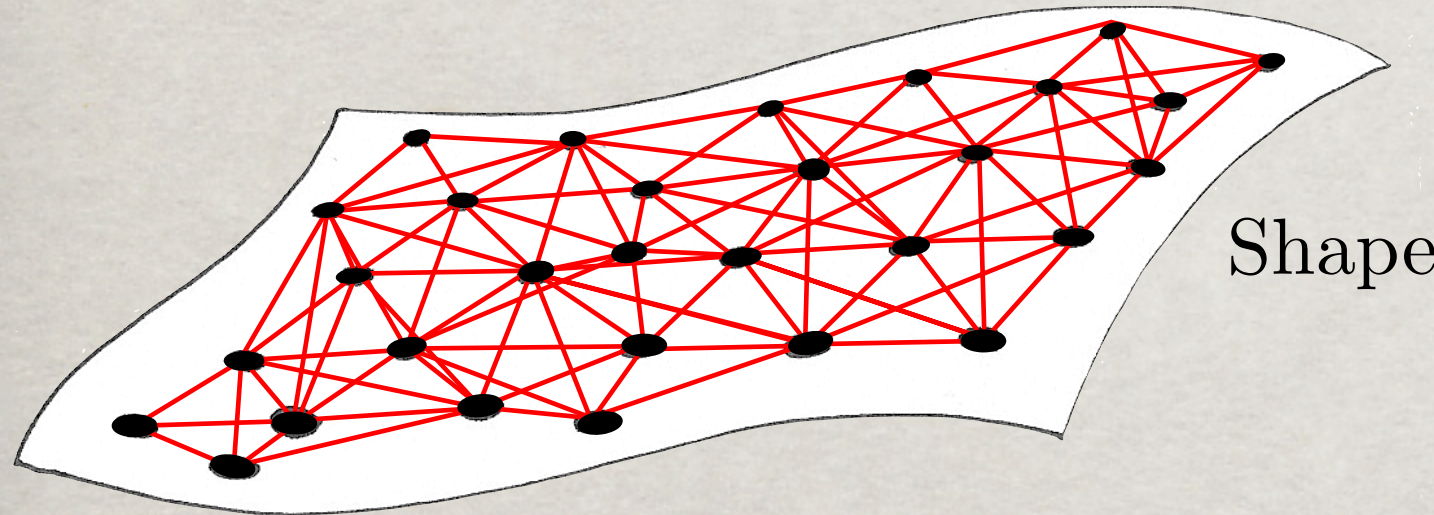
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SHAPE RECONSTRUCTION

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Point cloud P



OUTPUT

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\approx

for sampling conditions stated in

[AL 2010] when $d = d_\infty$

[ALS 2011] when $d = d_2$

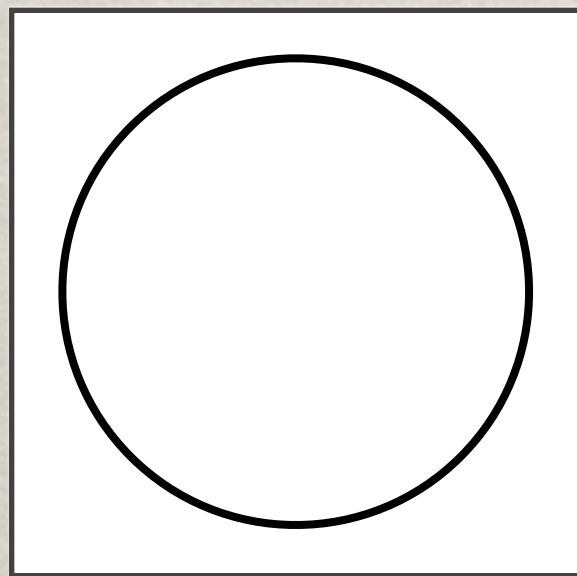
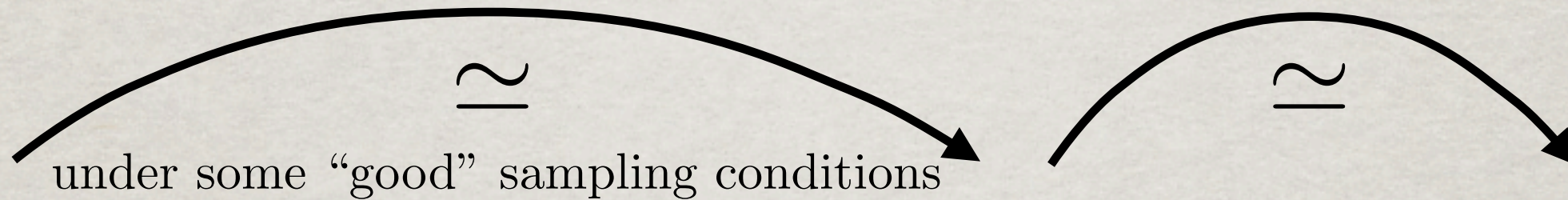
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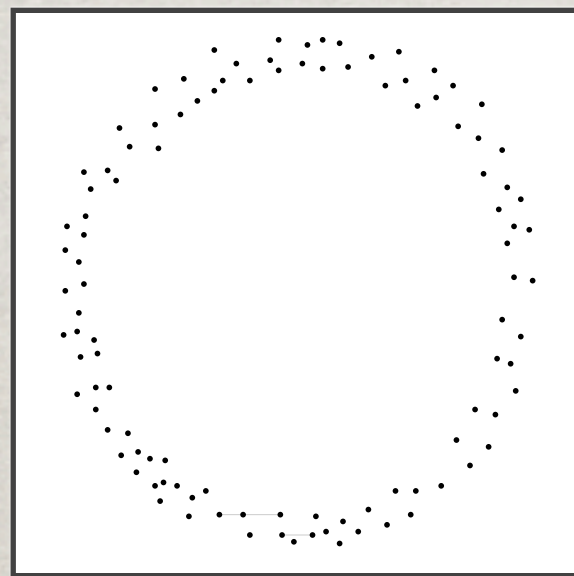
OVERVIEW

Part I:
Guarantees

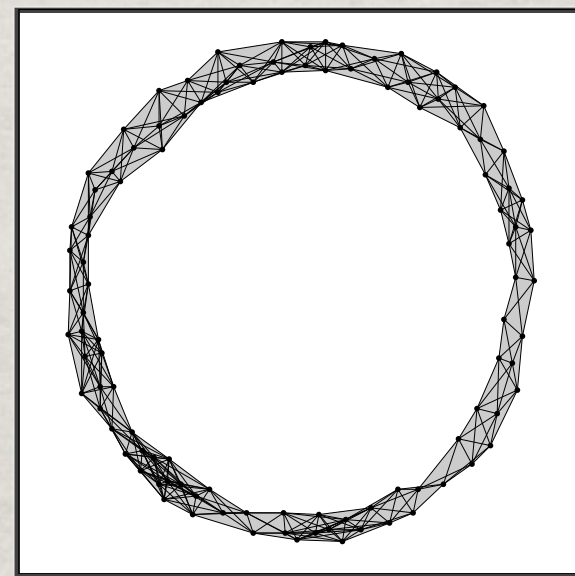
Part II:
Simplification



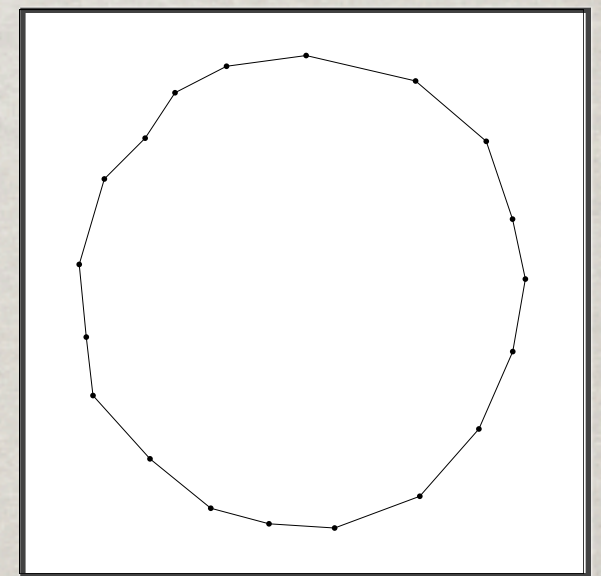
Shape



Point cloud



Flag complex

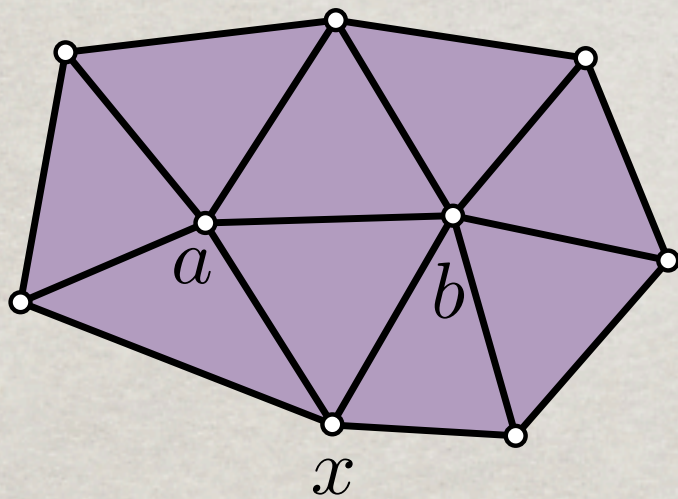


can be high dimensional !

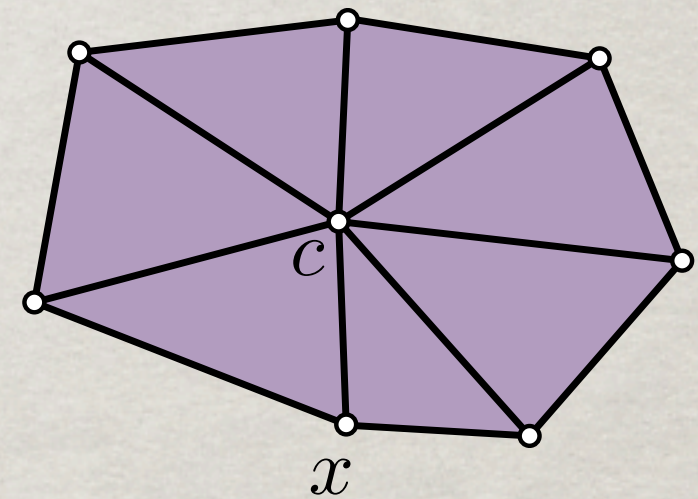
EDGE CONTRACTION

operation that identifies vertices a and b to vertex c

$$K \xrightarrow{ab \mapsto c} K'$$



$\{\dots, a, b, x, ab, ax, bx, abx, \dots\}$



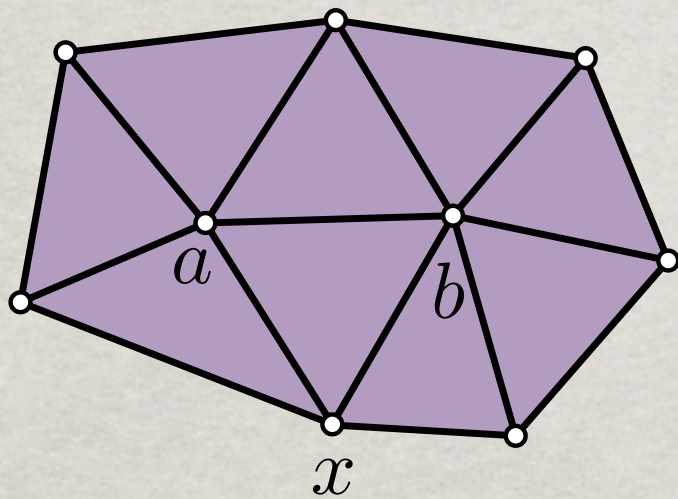
$\{\dots, c, x, cx, \dots\}$

EDGE CONTRACTION

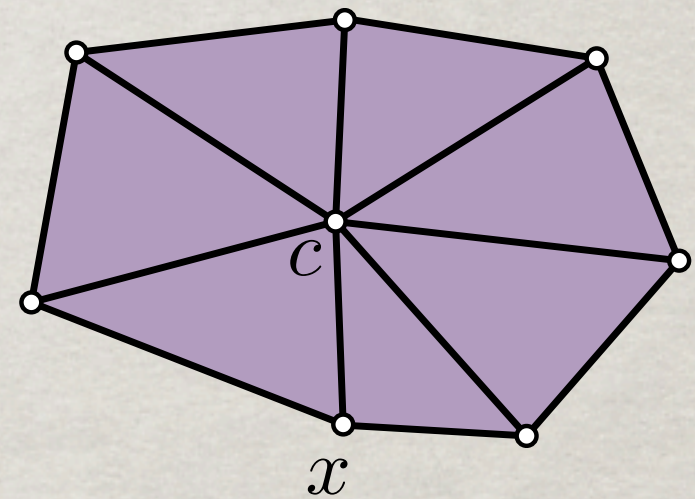
operation that identifies vertices a and b to vertex c

$$K \xrightarrow{ab \mapsto c} K' = \{f(\sigma) \mid \sigma \in K\}$$

$$f(v) = \begin{cases} c & \text{if } v \in \{a, b\} \\ v & \text{if } v \notin \{a, b\} \end{cases}$$



$\{\dots, a, b, x, ab, ax, bx, abx, \dots\}$



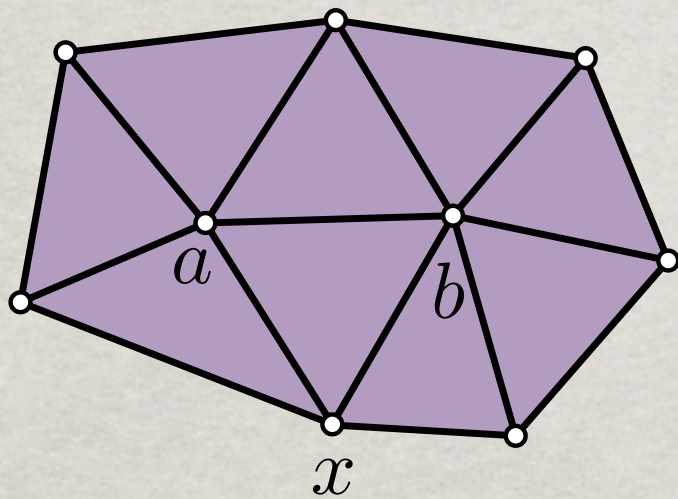
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EDGE CONTRACTION

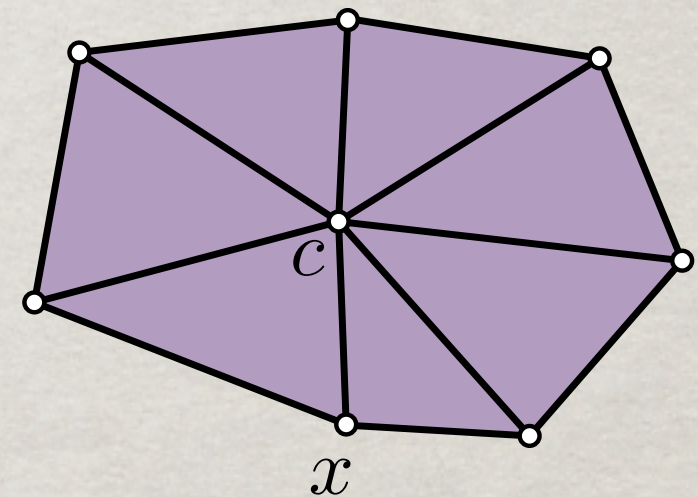
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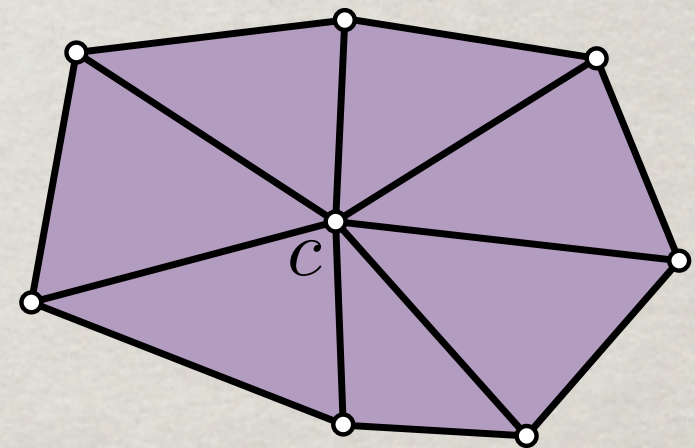
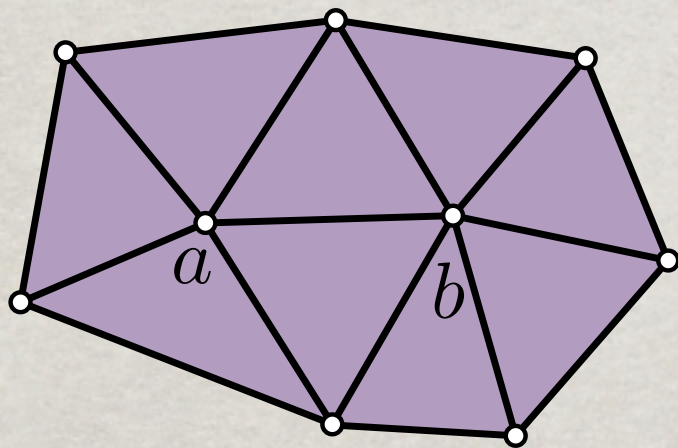


$\{\dots, c, x, cx, \dots\}$

- ✱ What if the result is not a flag complex?
- ✱ How to preserve homotopy type?

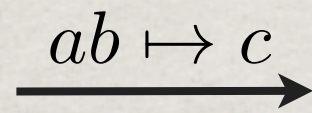
EDGE CONTRACTION

$$K = \text{Flag } K^{(1)} \xrightarrow{ab \mapsto c} K' = \text{Flag } K'^{(1)}$$



EDGE CONTRACTION

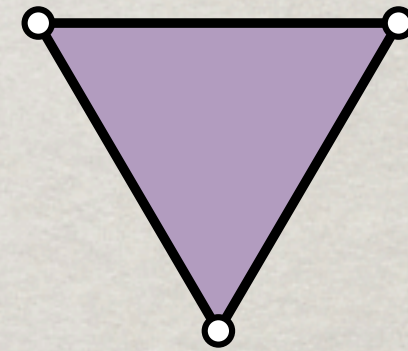
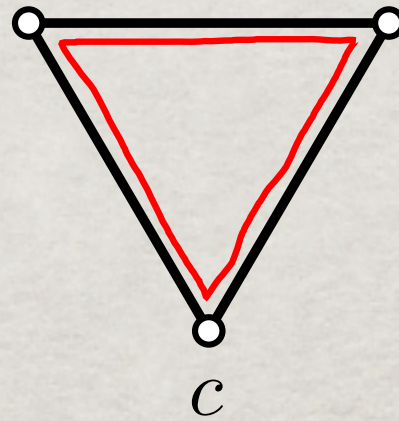
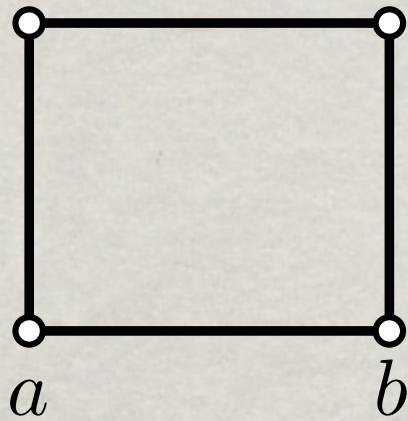
$$K = \text{Flag } K^{(1)}$$



$$K'$$

$$\neq$$

$$\text{Flag } K'^{(1)}$$



EDGE CONTRACTION

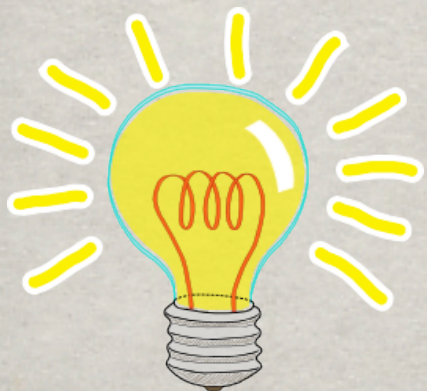
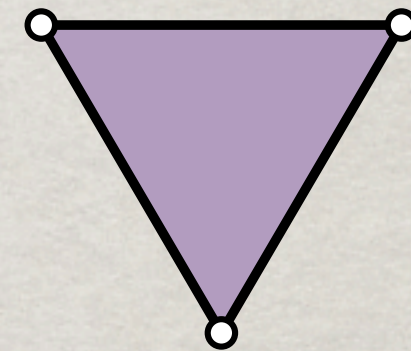
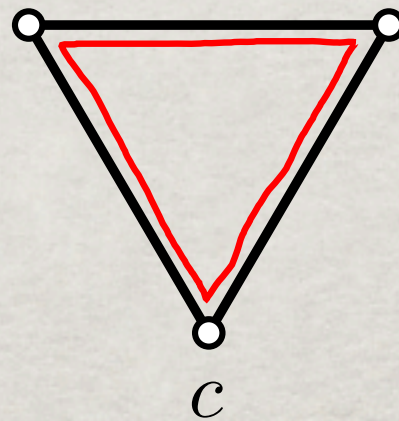
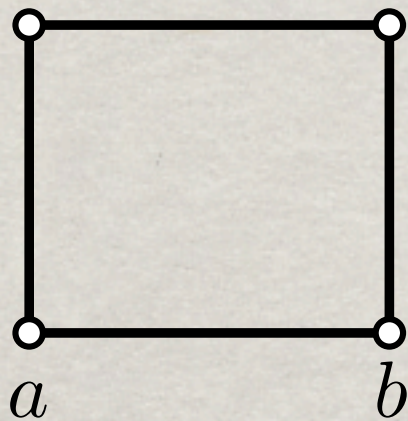
$$K = \text{Flag } K^{(1)}$$

$$\xrightarrow{ab \mapsto c}$$

$$K'$$

$$\neq$$

$$\text{Flag } K'^{(1)}$$



Encode a simplicial complex K by storing the pair:

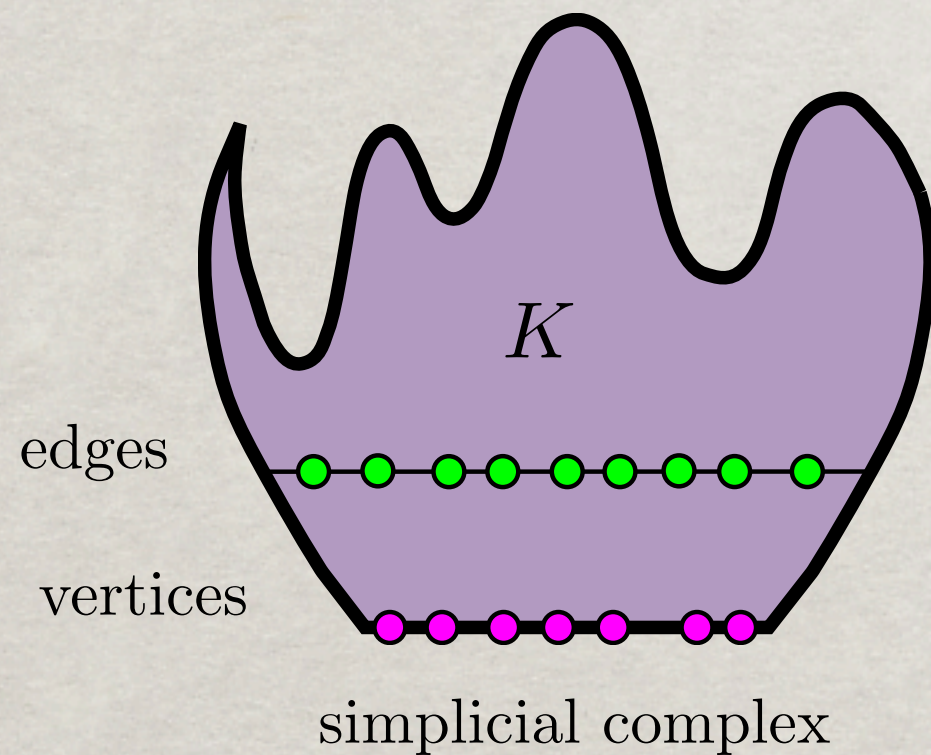
$$(K^{(1)}, \text{Blockers}(K))$$

vertices and edges

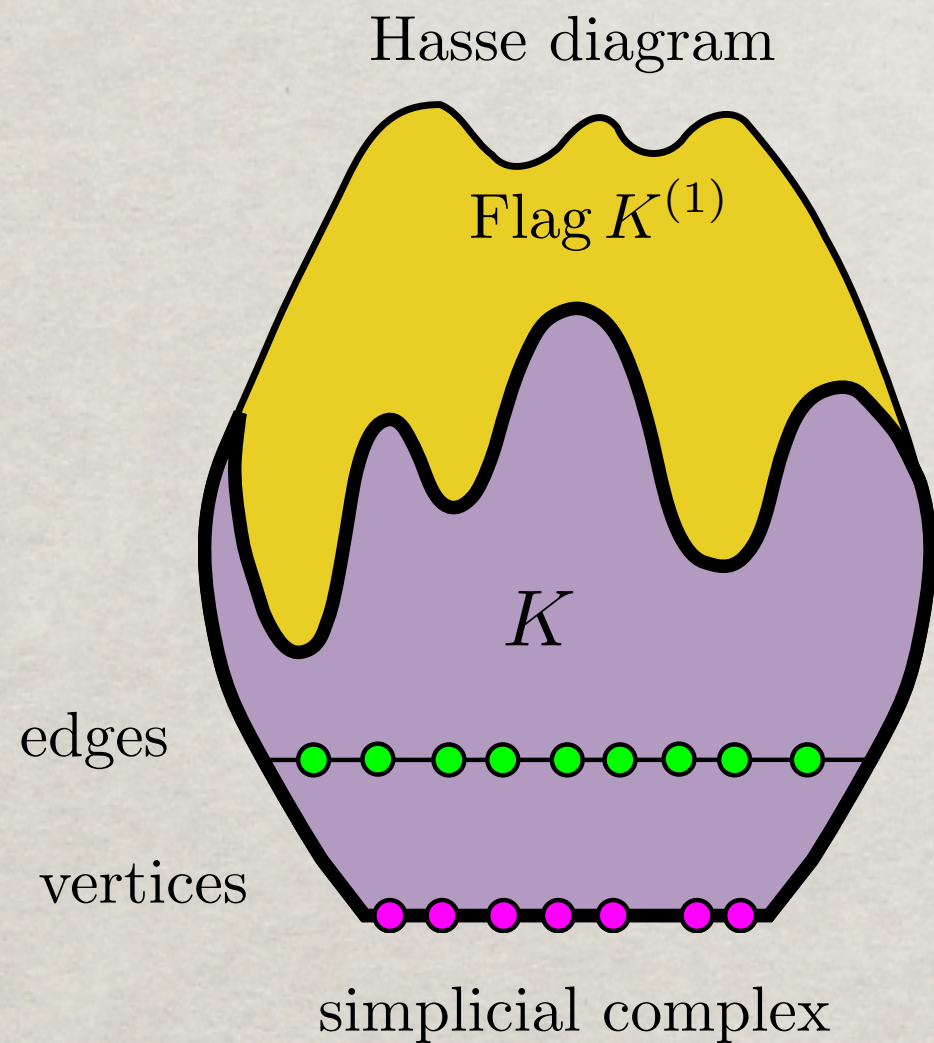
indicates how much
 K differs from $\text{Flag } K^{(1)}$

DATA STRUCTURE FOR SIMPLICIAL COMPLEXES

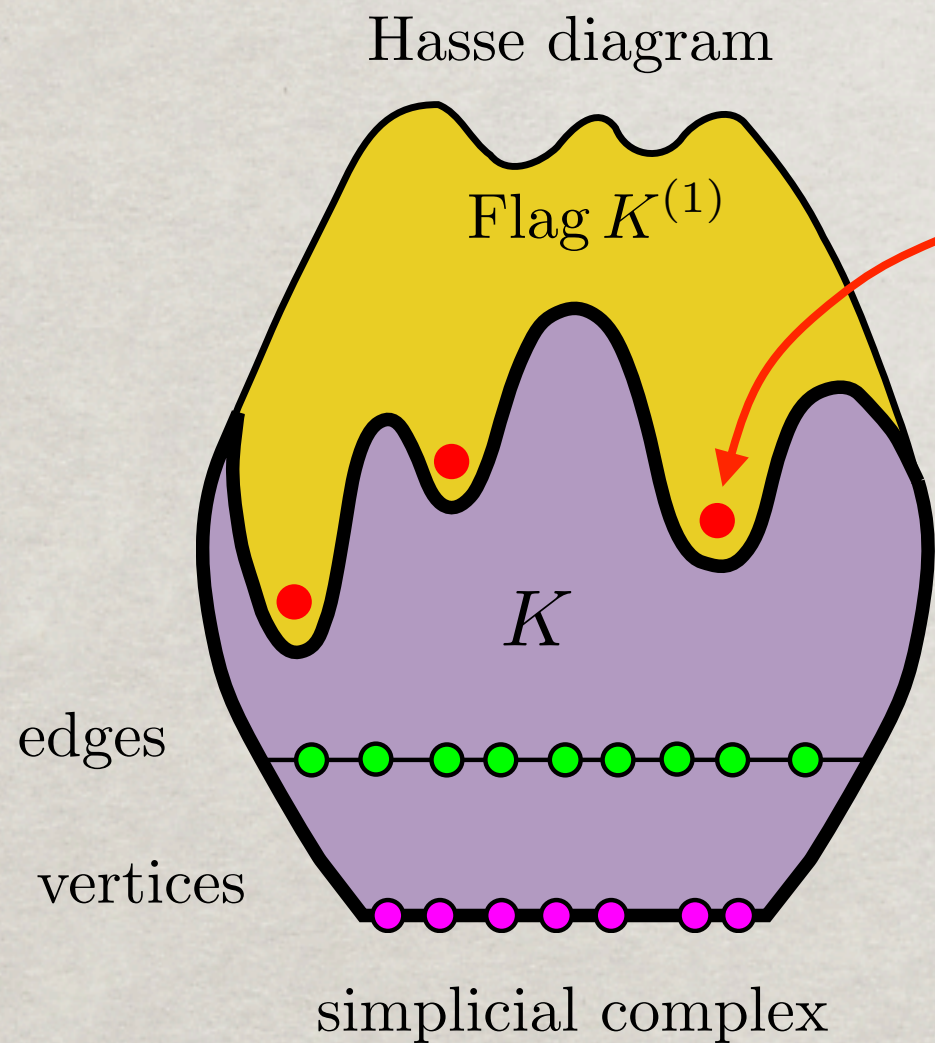
Hasse diagram



DATA STRUCTURE FOR SIMPLICIAL COMPLEXES

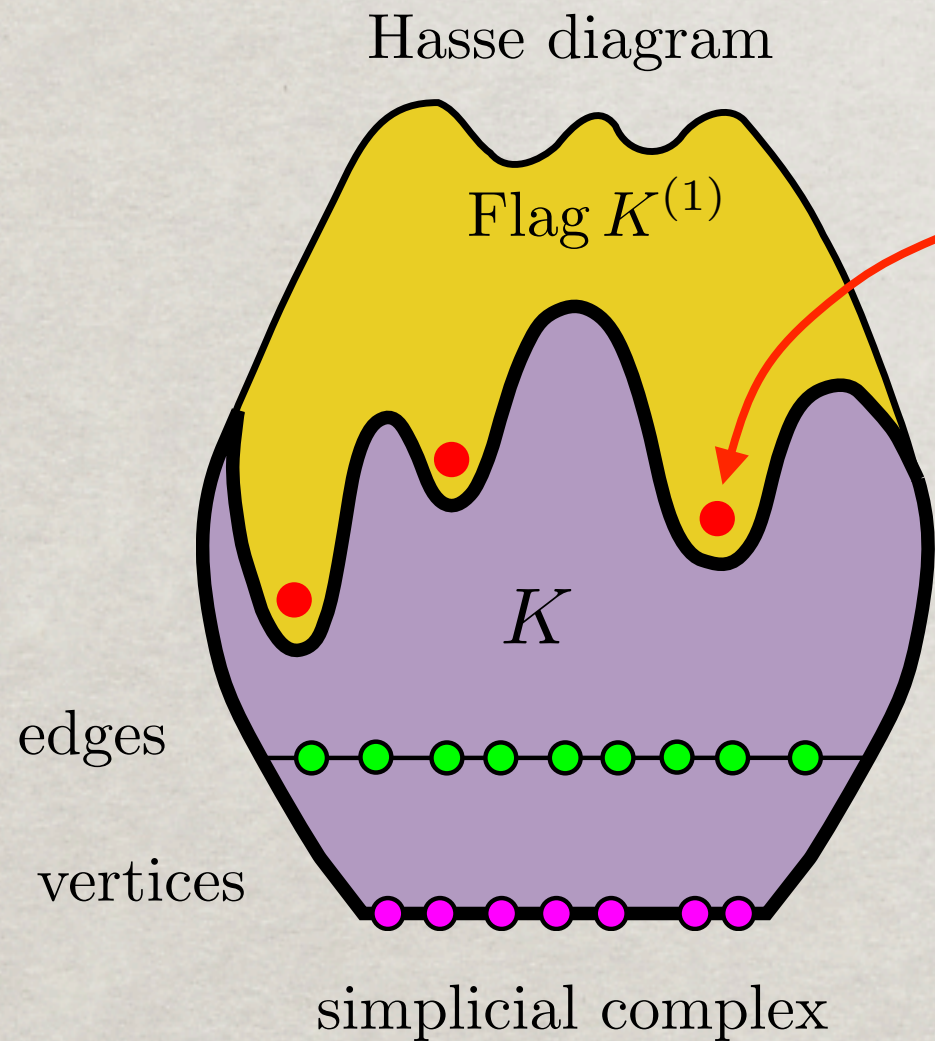


DATA STRUCTURE FOR SIMPLICIAL COMPLEXES



Blockers of K are inclusion-minimal simplices of $\text{Flag } K^{(1)} \setminus K$

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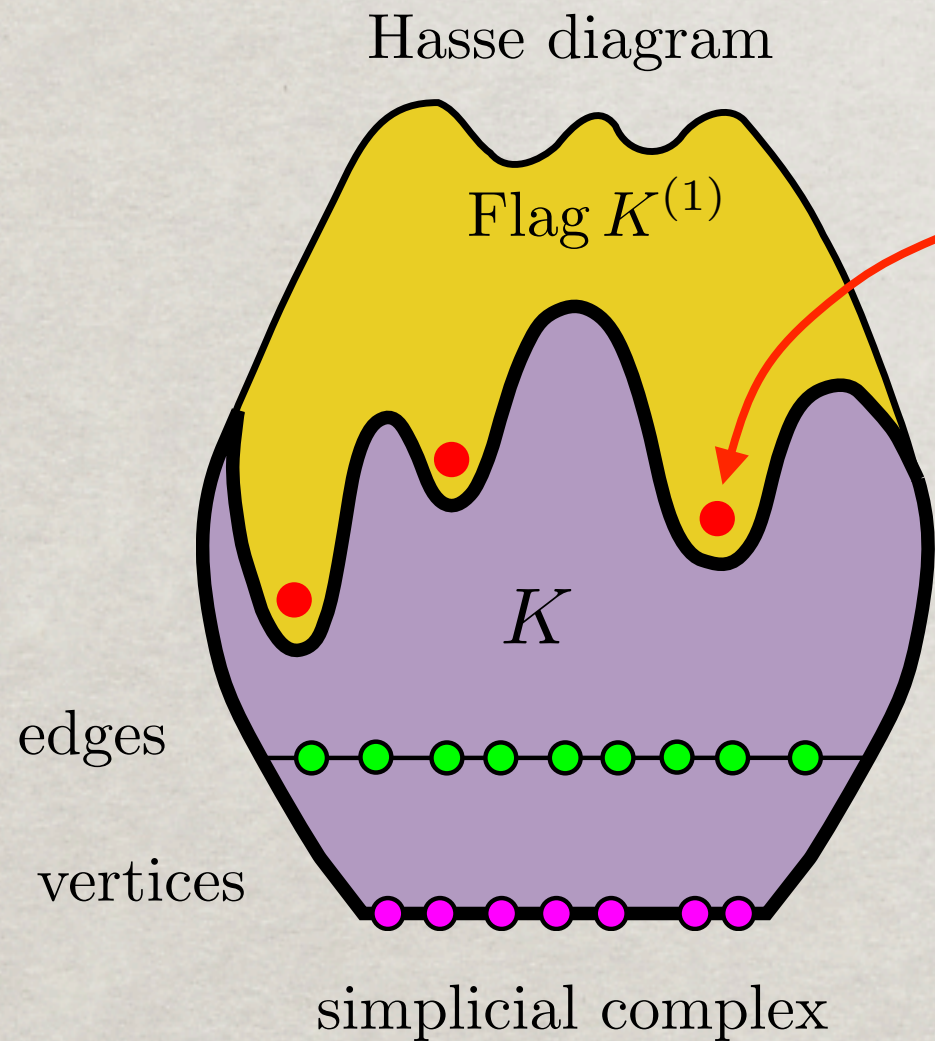
\iff

$\dim \sigma \geq 2$

$\sigma \notin K$

$\forall \tau \subsetneq \sigma, \tau \in K$

DATA STRUCTURE FOR SIMPLICIAL COMPLEXES



Blockers of K are inclusion-minimal simplices of $\text{Flag } K^{(1)} \setminus K$

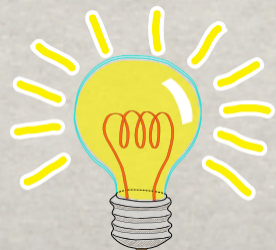
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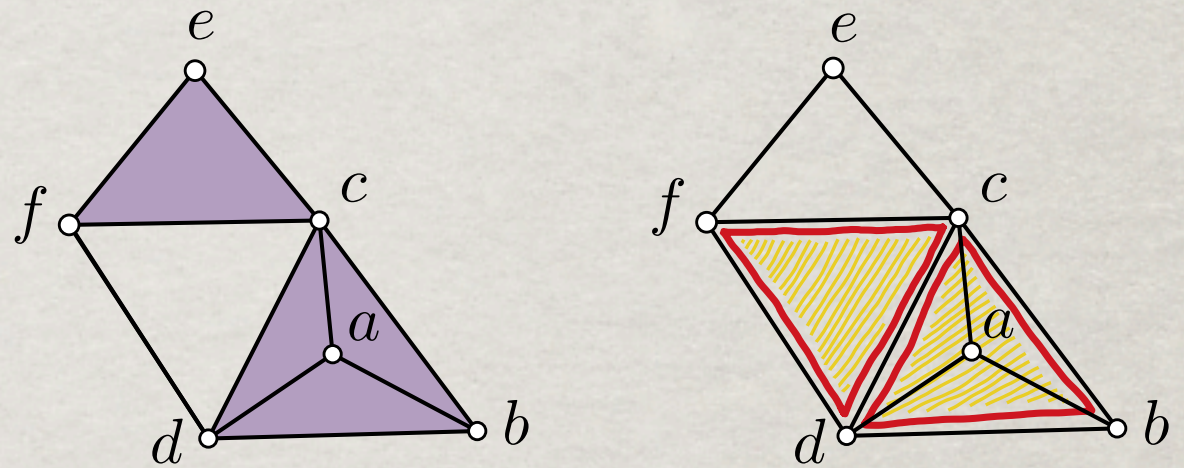
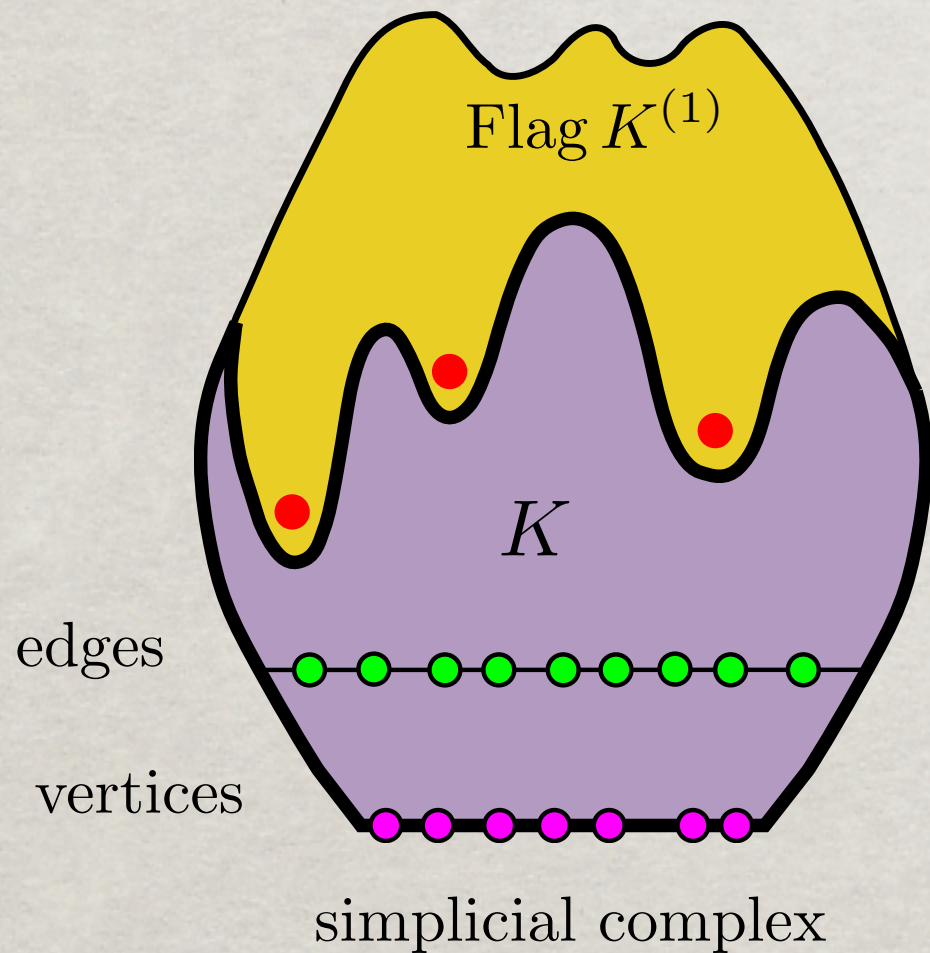
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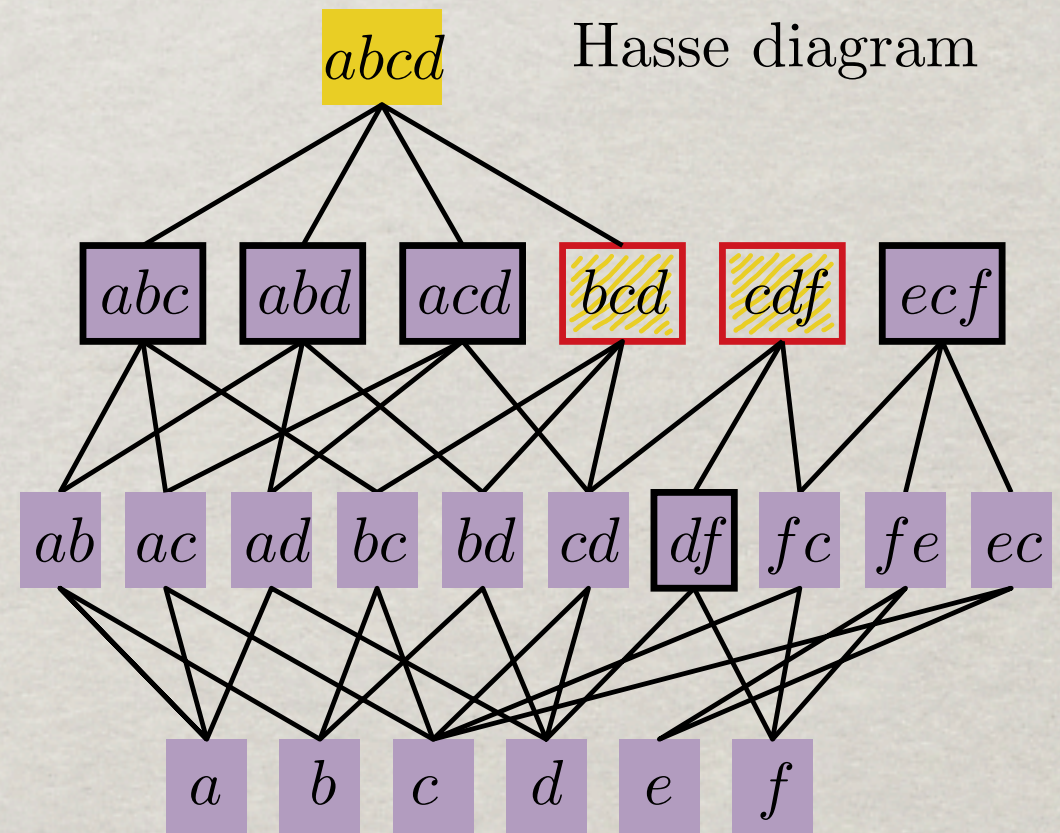
Encode a simplicial complex K by storing the pair:
 $(K^{(1)}, \text{Blockers}(K))$

A SMALL EXAMPLE

Hasse diagram



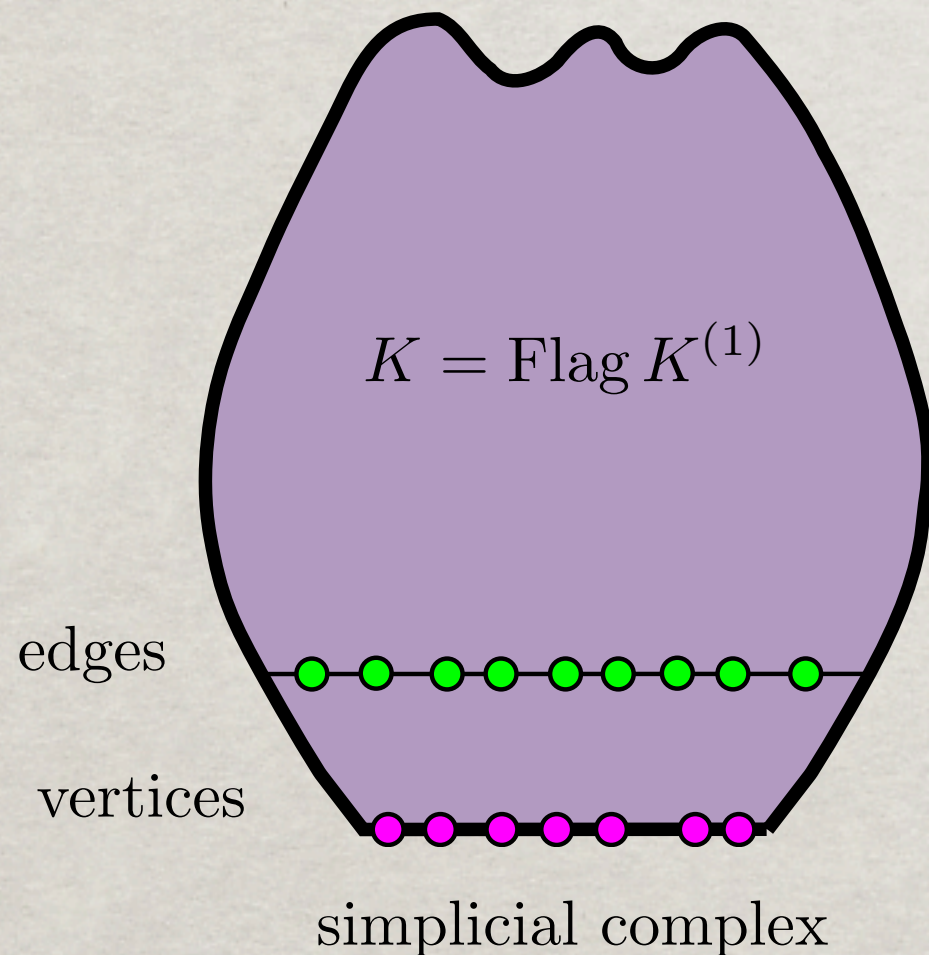
Hasse diagram



$$\text{Blockers}(K) = \{bcd, cdf\}$$

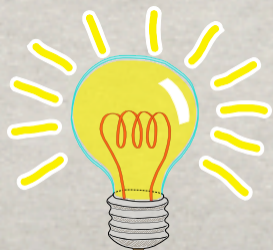
DATA STRUCTURE FOR SIMPLICIAL COMPLEXES

Hasse diagram



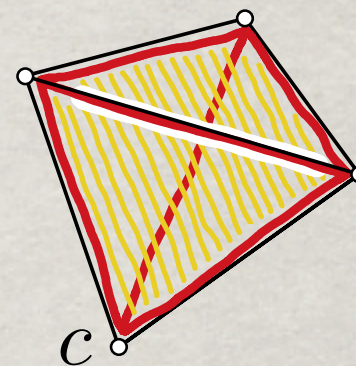
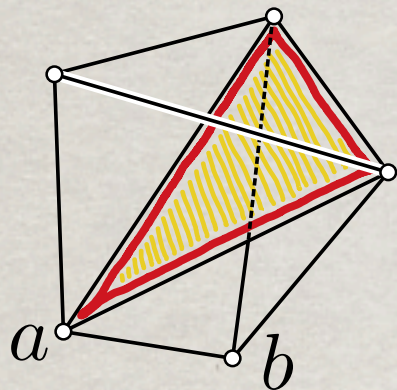
Blockers of K are inclusion-minimal simplices of $\text{Flag } K^{(1)} \setminus K$

If K is a flag complex
 $\text{Blockers}(K) = \emptyset$



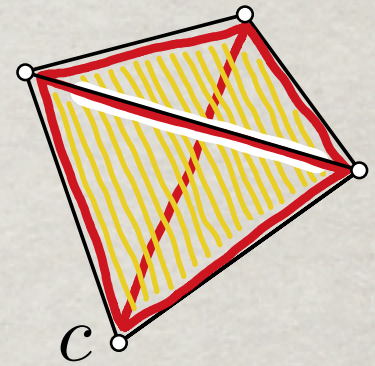
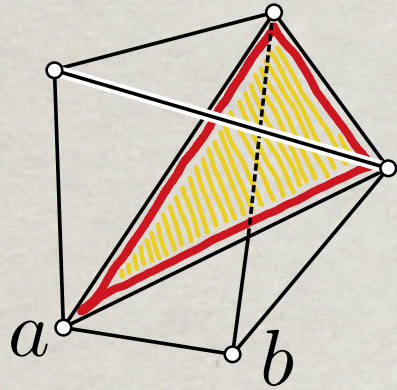
For a flag complex K , the pair reduces to:
 $(K^{(1)}, \emptyset)$

UPDATING DATA STRUCTURE



$$K = (K^{(1)}, \text{Blockers}(K)) \xrightarrow{ab \mapsto c} K' = (K'^{(1)}, \text{Blockers}(K'))$$

UPDATING DATA STRUCTURE



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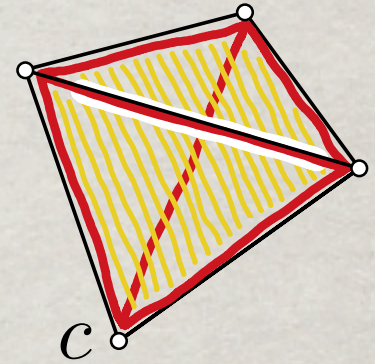
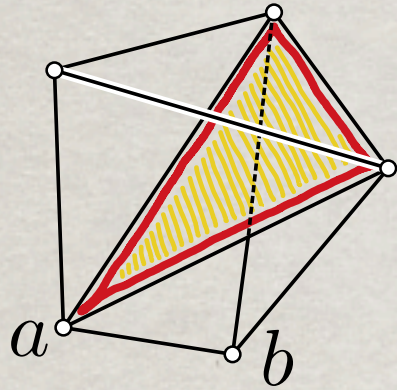
Lemma 1. $c\sigma \in \text{Blockers}(K')$ with $\sigma \subset \text{Vert}(K) \setminus \{a, b\}$ and $\dim \sigma \geq 1$ iff:

(i) $\sigma \in K$; for all $\tau \subsetneq \sigma$, $\tau \in \text{Lk}(a) \cup \text{Lk}(b)$;

(ii) $\sigma = \alpha\beta$ with $a\beta \in \text{Blockers}_0(K)$ and $b\alpha \in \text{Blockers}_0(K)$,

where $\text{Blockers}_0(K) = \text{Blockers}(K) \cup \text{complement of } K^{(1)}$

UPDATING DATA STRUCTURE



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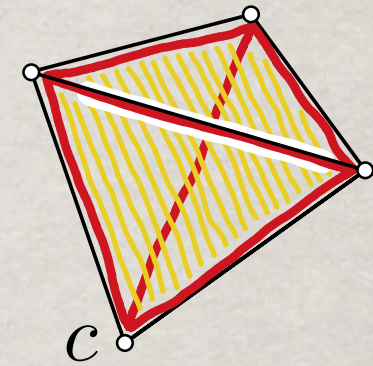
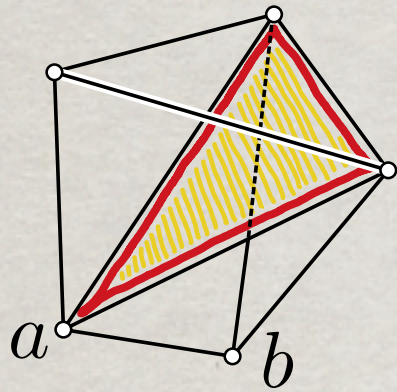
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$$\text{Lk } v = ((\text{Lk } v)^{(1)}, \text{Blockers}(\text{Lk } v))$$

UPDATING DATA STRUCTURE



$$K = (K^{(1)}, \text{Blockers}(K)) \xrightarrow{ab \mapsto c} K' = (K'^{(1)}, \text{Blockers}(K'))$$

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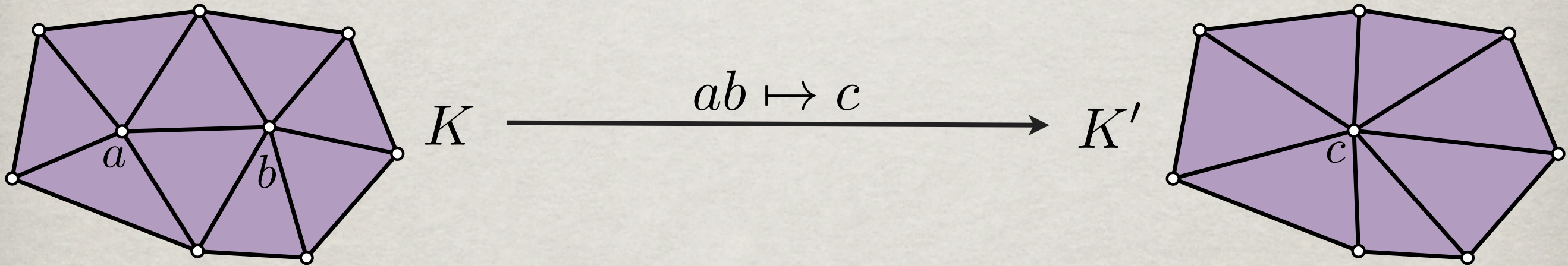


$$\text{Lk } v = ((\text{Lk } v)^{(1)}, \text{Blockers}(\text{Lk } v))$$

If no blockers “around” a and b , costs in $\tilde{O}(\#\text{neighbors}(a) \times \#\text{neighbors}(b))$

EDGE CONTRACTION

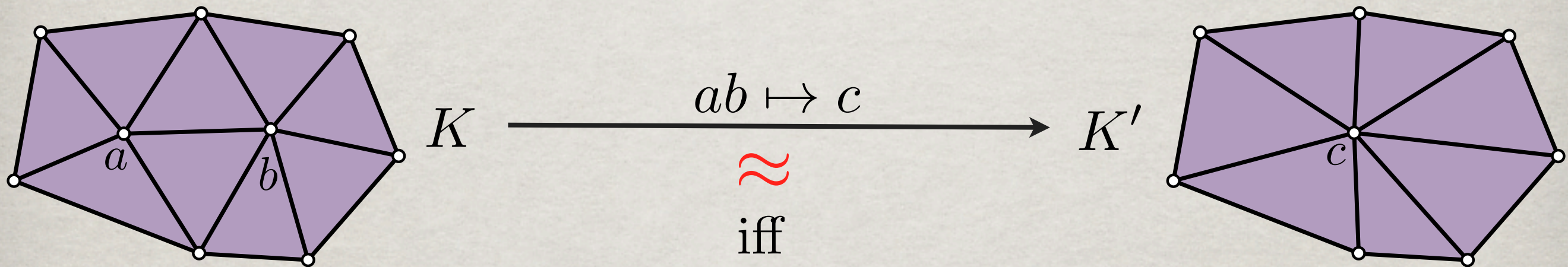
How to preserve homotopy type?



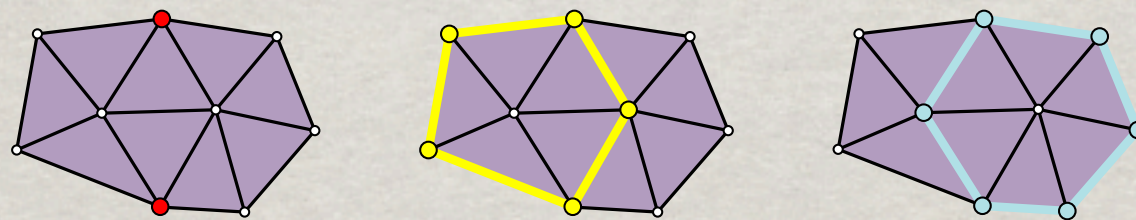
EDGE CONTRACTION

[Dey, Edelsbrunner, Guha & Nekhayev 1999]

If K triangulates a 2- or 3-manifold

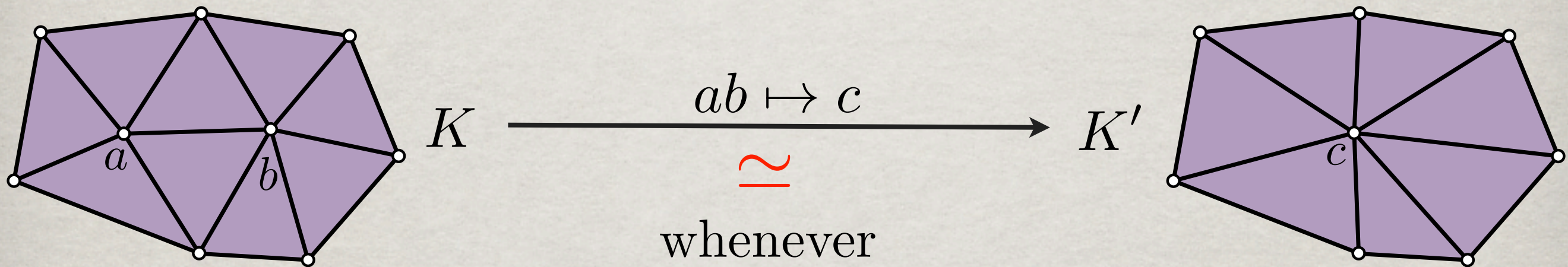


$$\text{Lk } ab = \text{Lk } a \cap \text{Lk } b$$

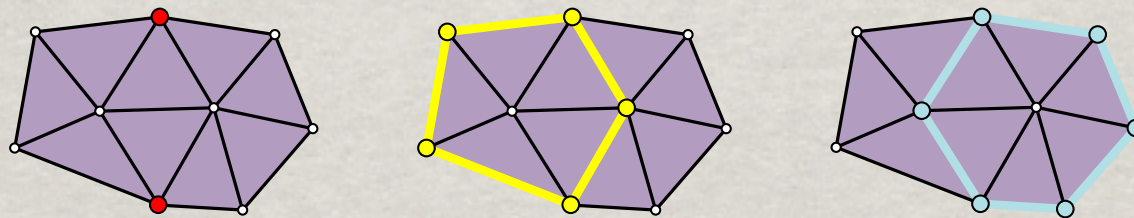


EDGE CONTRACTION

For arbitrary simplicial complexes, we established that:

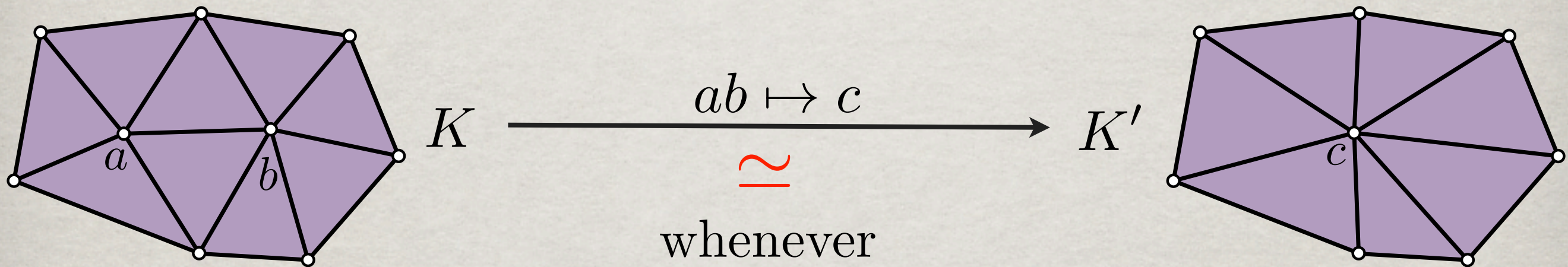


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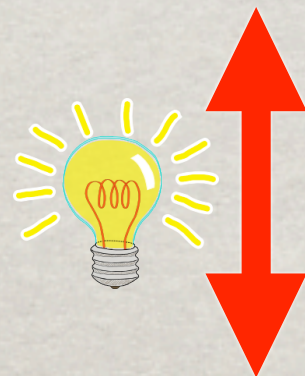


EDGE CONTRACTION

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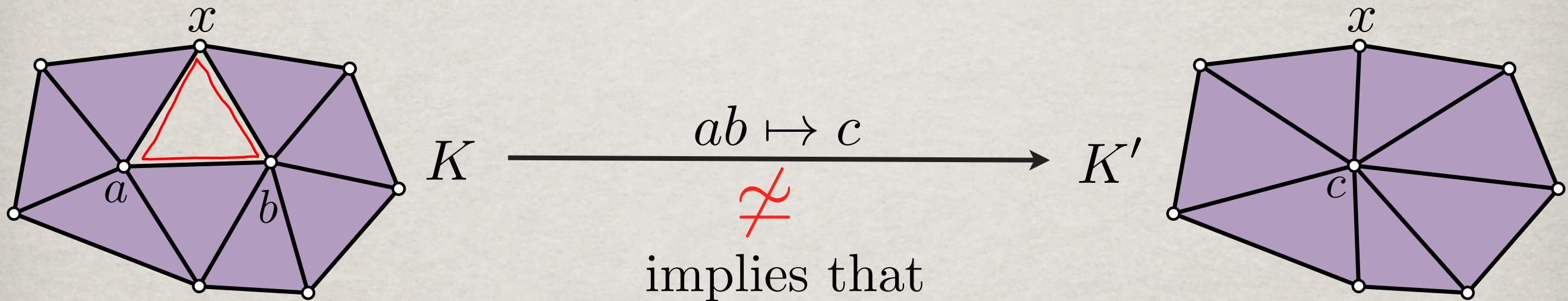
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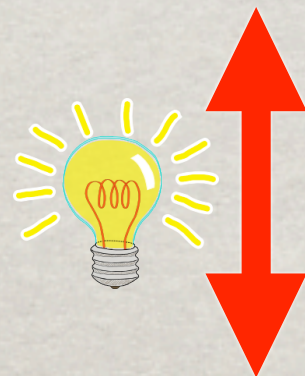
No blocker of K contains ab

EDGE CONTRACTION

For arbitrary simplicial complexes, we established that:



$$\text{Lk } ab \neq \text{Lk } a \cap \text{Lk } b$$



\exists a blocker of K containing ab

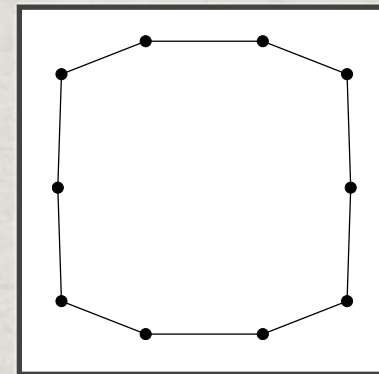
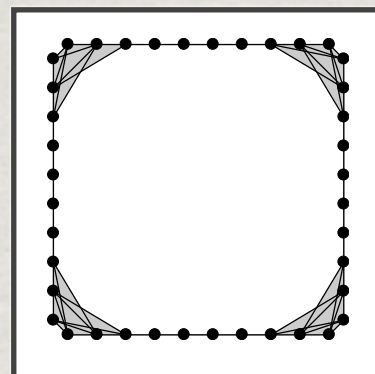
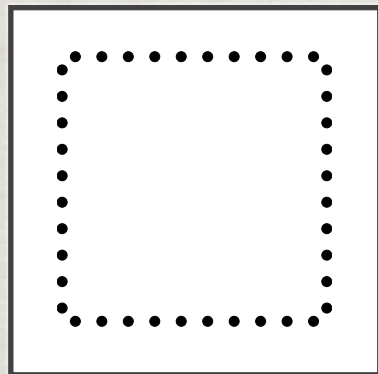
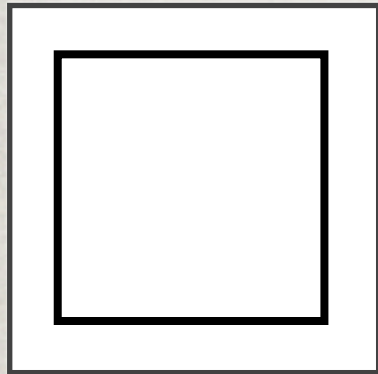
EXPERIMENTS

We keep contracting shortest edge with no blocker through it

$$C_d = \partial[-1, 1]^{d+1}$$

Point cloud

Flag complex

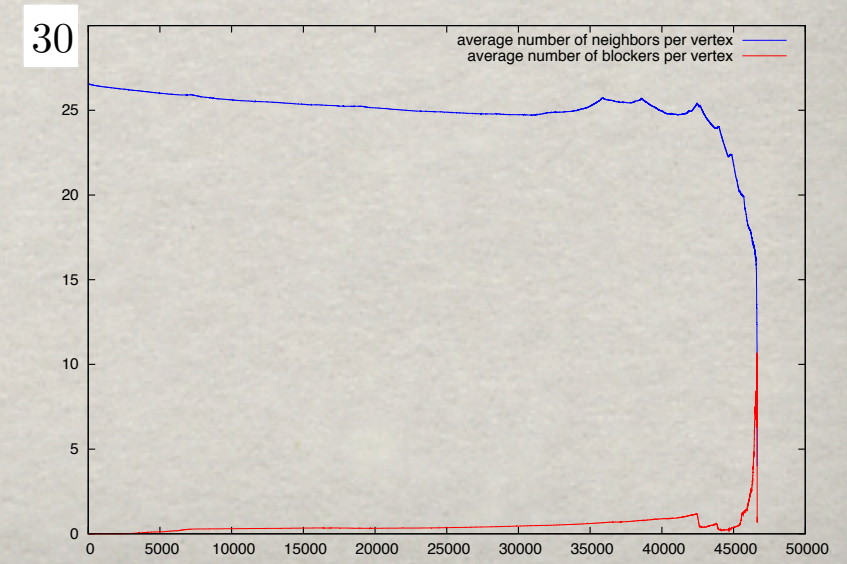
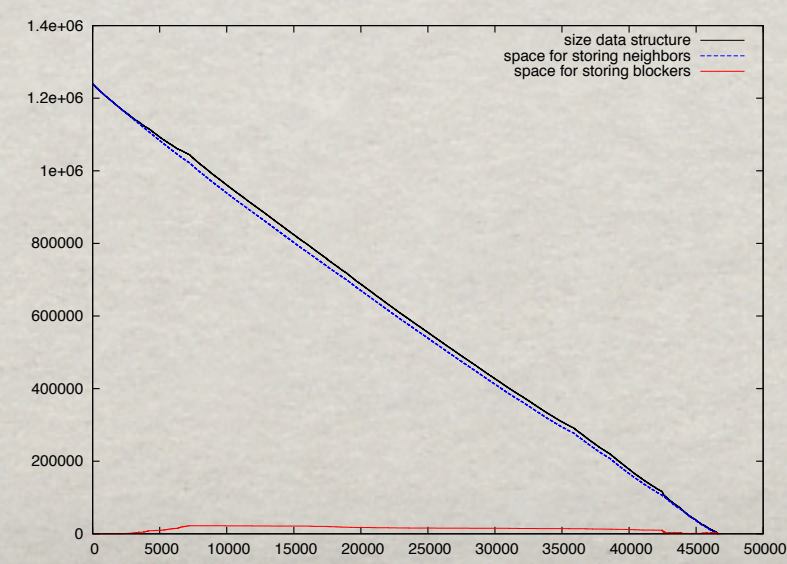
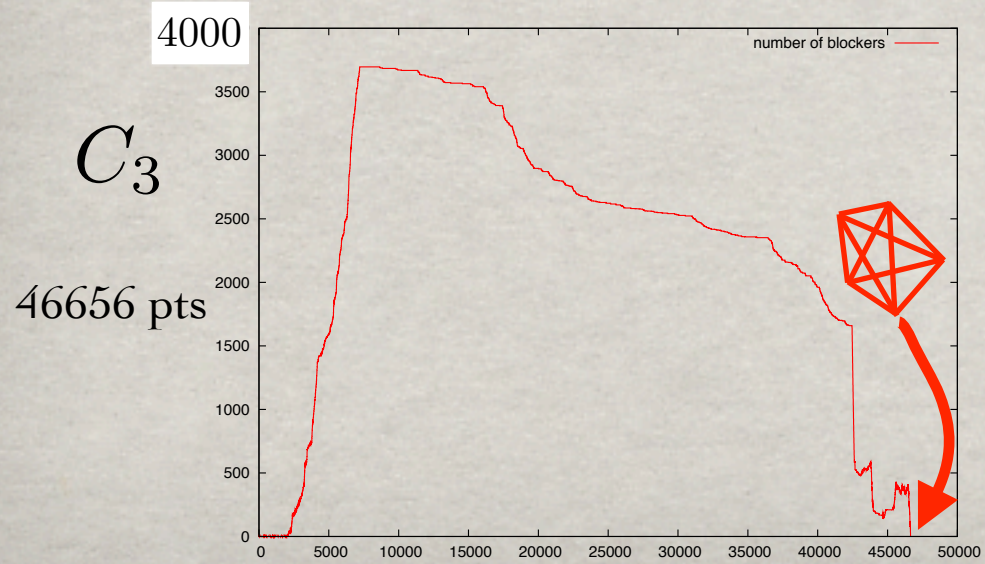
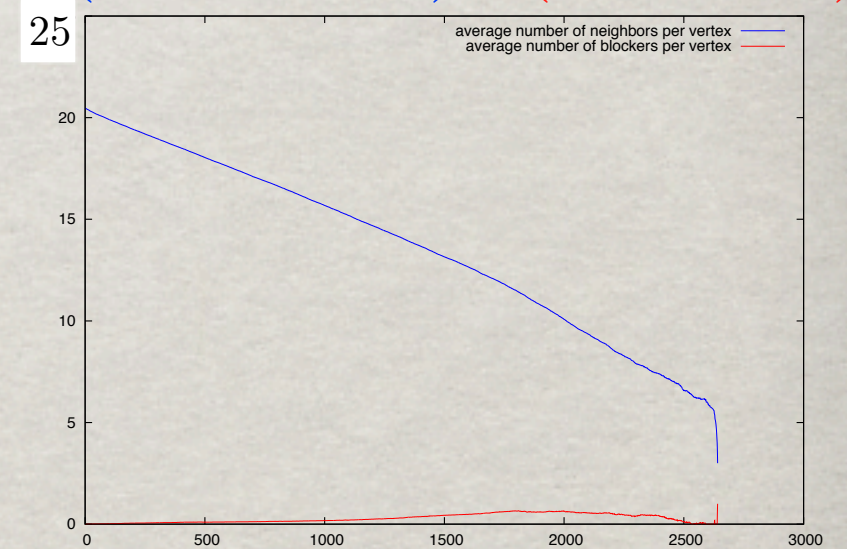
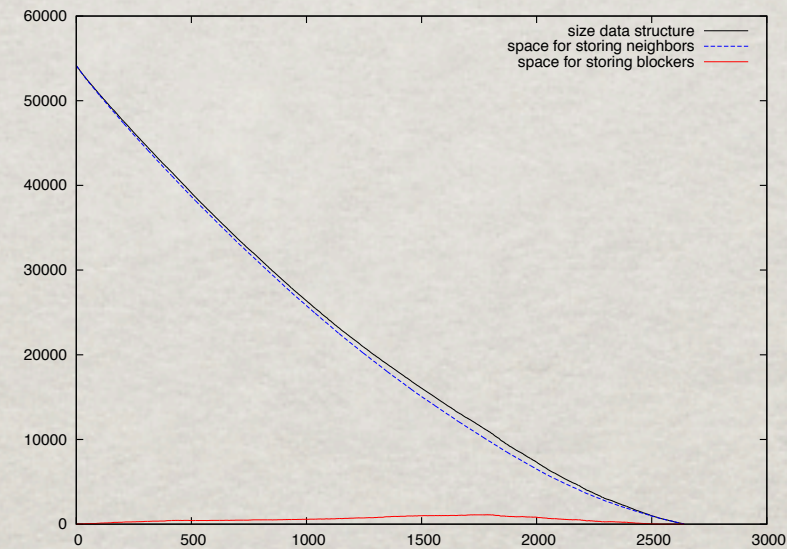
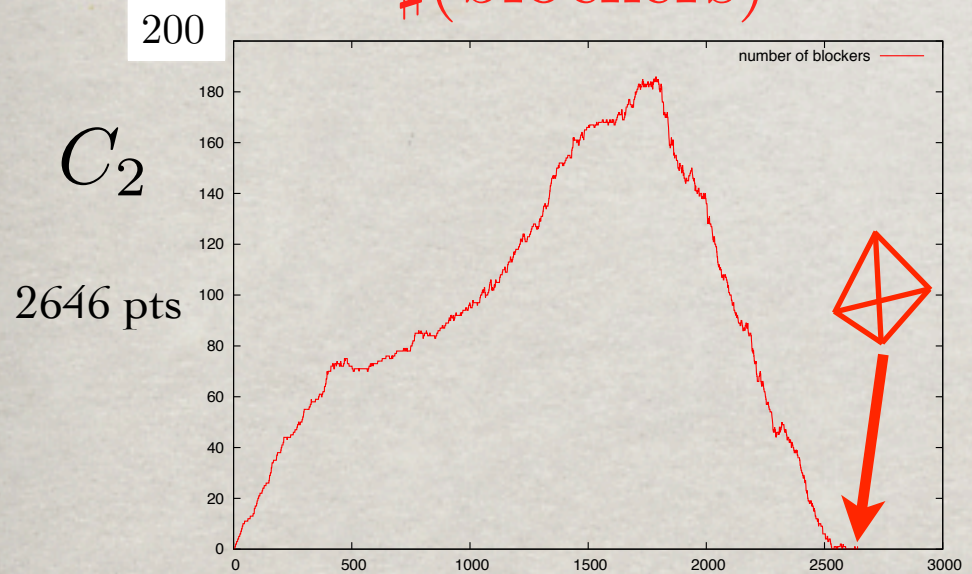


\approx

#(blockers)

size of data structure

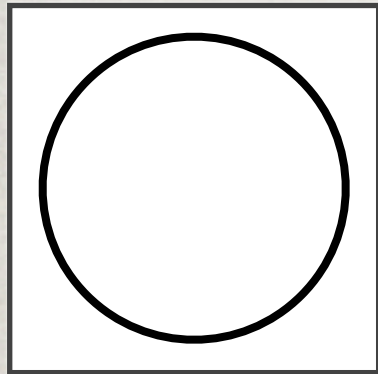
E(neighbors) E(blockers)



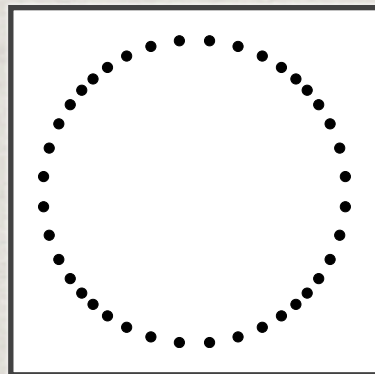
EXPERIMENTS

We keep contracting shortest edge with no blocker through it

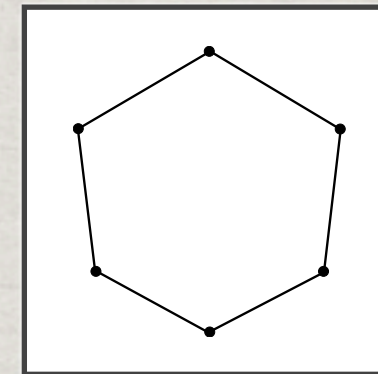
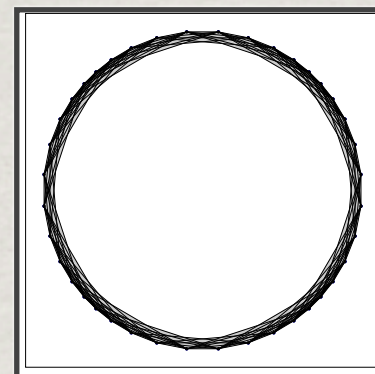
$S_d = d$ -sphere



Point cloud



Flag complex

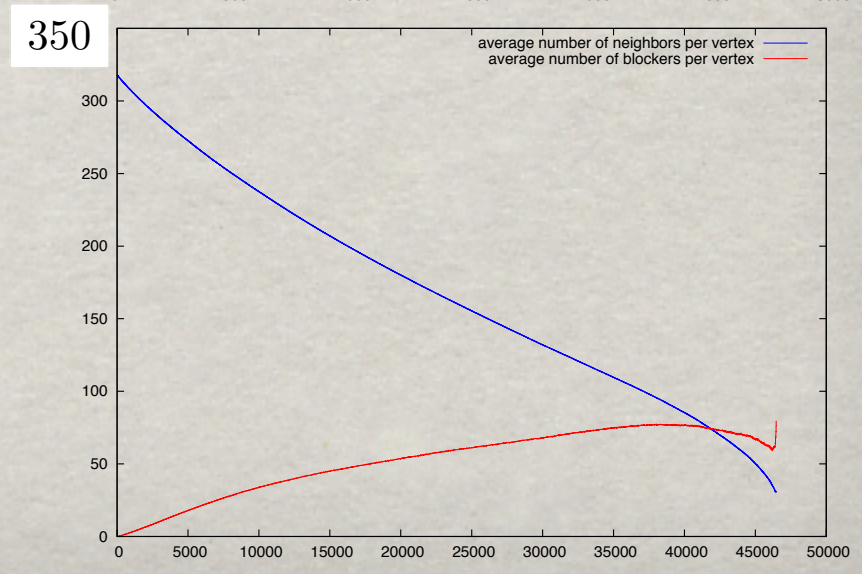
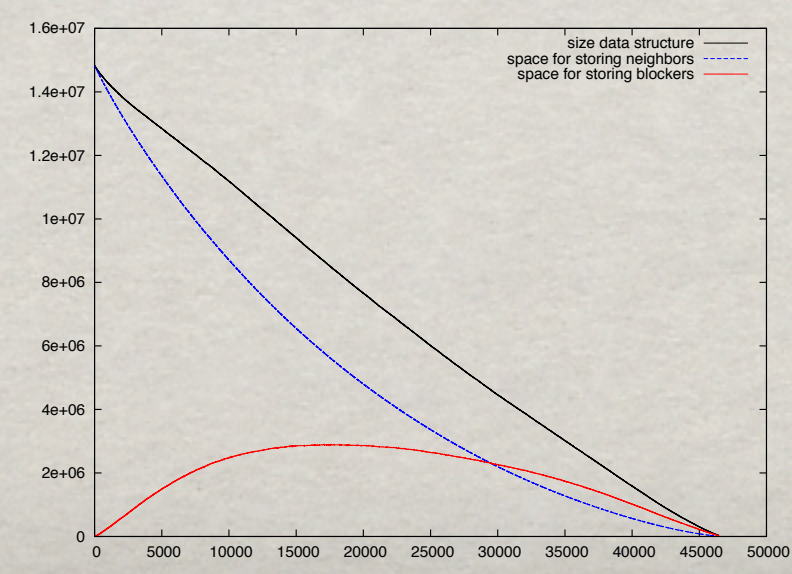
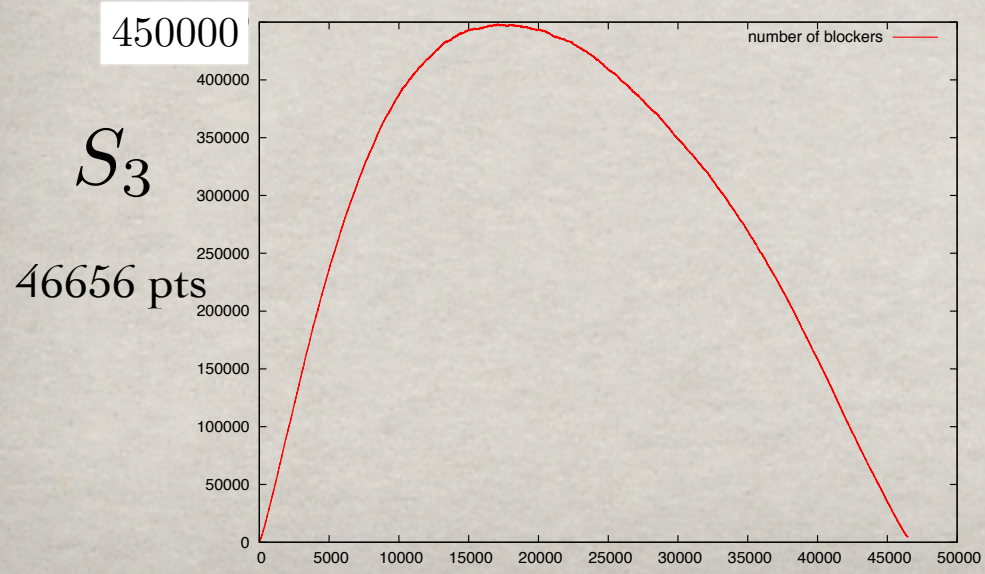
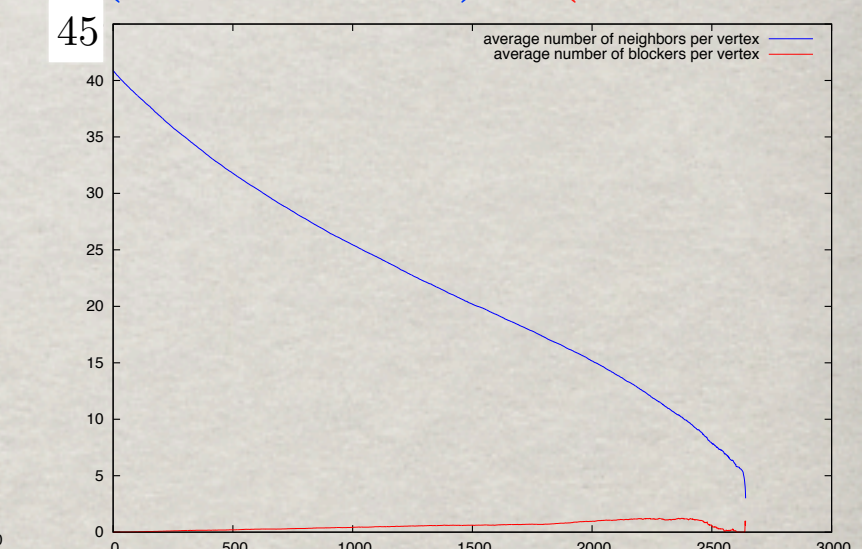
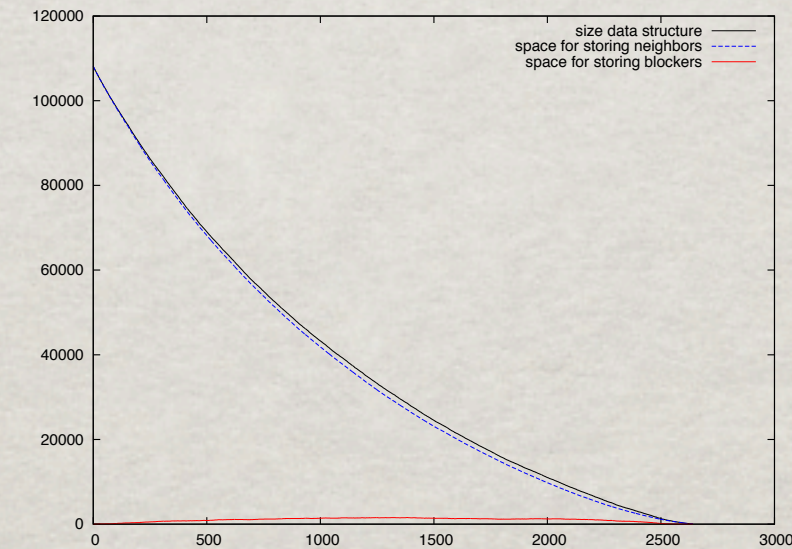
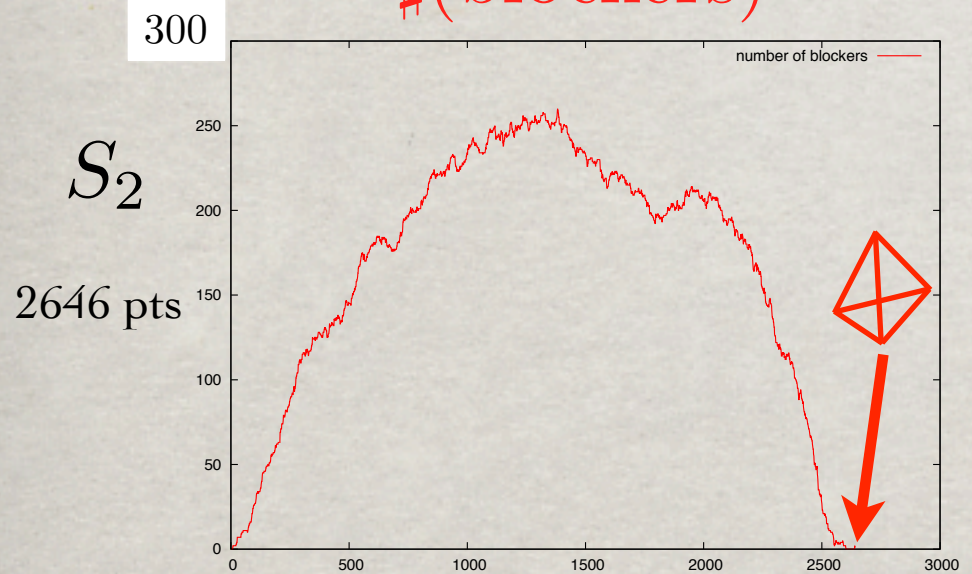


~

#(blockers)

size of data structure

$E(\text{neighbors})$ $E(\text{blockers})$

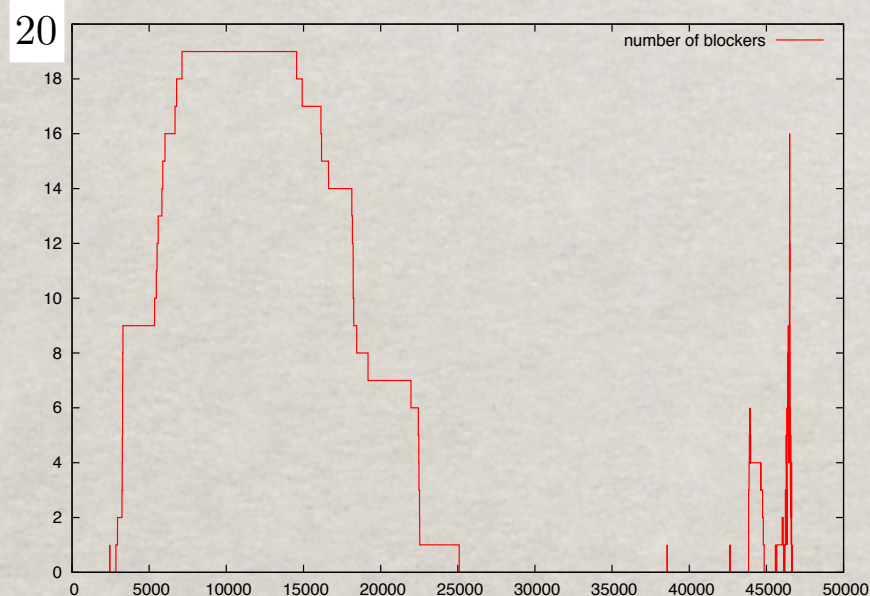


EXPERIMENTS

$$K = (G, B) \xrightarrow{\text{extended anticollapse}} K' = (G, B \setminus \{b\})$$

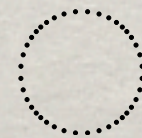
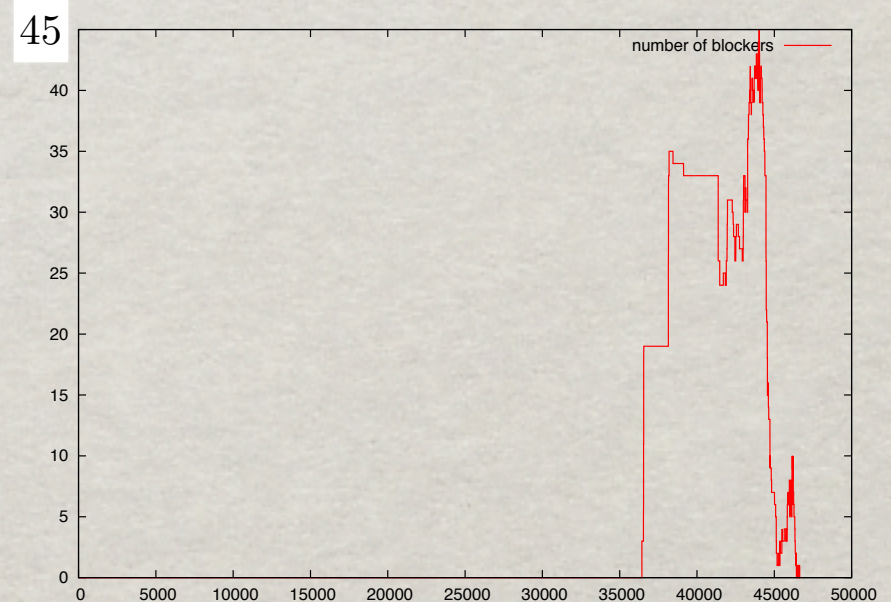
\approx
whenever

$\text{Lk}_{K'} b$ is a cone



C3

46656 pts



S3

46656 pts

CONCLUSION

- ✱ Promising data structure for encoding high dimensional simplicial complexes.
- ✱ Need further investigations:
 - ✱ study other strategies to prioritize contractions besides edge length
 - ✱ study other simplification operations besides edge contractions (collapses, ...)

